

Any Monotone Property of 3-uniform Hypergraphs is Weakly Evasive

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Abstract. For a Boolean function f , let $D(f)$ denote its deterministic decision tree complexity, i.e., minimum number of (adaptive) queries required in worst case in order to determine f . In a classic paper, Rivest and Vuillemin [19] show that any non-constant monotone property $\mathcal{P} : \{0, 1\}^{\binom{n}{2}} \rightarrow \{0, 1\}$ of n -vertex graphs has $D(\mathcal{P}) = \Omega(n^2)$.

We extend their result to 3-uniform hypergraphs. In particular, we show that any non-constant monotone property $\mathcal{P} : \{0, 1\}^{\binom{n}{3}} \rightarrow \{0, 1\}$ of n -vertex 3-uniform hypergraphs has $D(\mathcal{P}) = \Omega(n^3)$.

Our proof combines the combinatorial approach of Rivest and Vuillemin with the topological approach of Kahn, Saks, and Sturtevant. Interestingly, our proof makes use of Vinogradov's Theorem (weak Goldbach Conjecture), inspired by its recent use by Babai et. al. [1] in the context of the topological approach. Our work leaves the generalization to k -uniform hypergraphs as an intriguing open question.

Keywords: decision tree complexity; hypergraph properties; evasiveness; query complexity; group actions

1 Introduction

The *decision tree model* aka *query model* [3], perhaps due to its simplicity and fundamental nature, has been extensively studied over decades; yet there remain some outstanding open questions about it.

Fix a Boolean function $f : \{0, 1\}^n \rightarrow \{0, 1\}$. A deterministic decision tree D_f for f takes $x = (x_1, \dots, x_n)$ as an input and determines the value of $f(x_1, \dots, x_n)$

* Part of this work was done while the author was visiting the Centre for Quantum Technologies, National University of Singapore. This work was supported in part by the National Natural Science Foundation of China Grant 61170062, 61222202.

** Research at the Centre for Quantum Technologies is funded by the Singapore Ministry of Education and the National Research Foundation.

using queries of the form “is $x_i = 1$?”. Let $C(D_f, x)$ denote the cost of the computation, that is the number of queries made by D_f on input x . The *deterministic decision tree complexity* of f is defined as $D(f) = \min_{D_f} \max_x C(D_f, x)$.

The function f is called *evasive* if $D(f) = n$, i.e., one must query all the variables in worst case in order to determine the value of the function.

1.1 The Anderaa-Rosenberg-Karp Conjecture

A Boolean function f is said to be monotone (increasing) if for any $x \leq y$ we have $f(x) \leq f(y)$, where $x \leq y$ iff for all $i : x_i \leq y_i$. A property of n -vertex graphs is a Boolean function $\mathcal{P} : \{0, 1\}^{\binom{n}{2}} \rightarrow \{0, 1\}$ whose variables are identified with the $\binom{n}{2}$ potential edges of n -vertex graphs and the function \mathcal{P} is invariant under relabeling of the vertices. $\mathcal{P}(G) = 0$ means that the graph G satisfies the property. A natural theme in the study of decision tree complexity is to exploit the structure within f to prove strong lower bounds on its query complexity. A classic example is the following conjecture attributed to Anderaa, Rosenberg, and Karp, asserting the *evasiveness* of monotone graph properties:

Conjecture 1 (ARK Conjecture). (cf. [8]) Every non-trivial monotone graph property is evasive.

Some natural examples of monotone graph properties are: connectedness, planarity, 3-colorability, containment of a fixed subgraph etc.

Since its origin around 1975, the ARK Conjecture has caught the imagination of generations of researchers resulting in beautiful mathematical ideas; yet - to this date - remains unsolved. A major breakthrough on ARK Conjecture was obtained by Kahn, Saks, and Sturtevant [8] via their novel *topological approach*. They settled the conjecture when the number of vertices of the graphs is a power of prime number. The topological approach subsequently turned out useful for solving some other variants and special cases of the conjecture. For example: Yao confirms the variant of the conjecture for monotone properties of bipartite graphs [25]. More recently, building on Chakraborty, Khot, and Shi's work [4], Babai et. al. [1] show that under some well-known conjectures in number theory, *forbidden subgraph* property - containment of a fixed subgraph in the graph - is evasive. We refer the readers to the lecture notes [12] by Lovász and Young for a nice exposition of the works around this topic.

1.2 The Evasiveness Conjecture

The key feature of monotone graph properties is that they are sufficiently *symmetric*. In particular, they are *transitive* Boolean functions, i.e., there is a group acting transitively on the set of variables under which the function remains invariant. A natural question was raised: how much *symmetry* is necessary in order to guarantee the evasiveness? The following generalization (cf. [13]) of ARK Conjecture asserts that only transitivity suffices.

Conjecture 2 (Evasiveness Conjecture (EC)). If f is a non-trivial monotone transitive Boolean function, then f is evasive.

Rivest and Vuillemin [19] confirm the above conjecture when the number of variables is a power of prime number. The general case remains widely open.

1.3 The Weak Evasiveness Conjecture

Recently Kulkarni [7] proposes to investigate the following:

Conjecture 3 (Weak Evasiveness Conjecture). If $\{f_n\}$ is a sequence of non-trivial monotone transitive Boolean functions then for every $\epsilon > 0$

$$D(f_n) \geq n^{1-\epsilon}.$$

The best known lower bound in this context is $D(f) \geq R(f) \geq n^{2/3}$, which follows from the work of O'Donnell et. al. [18]. It turns out that [7] the above conjecture is equivalent to the EC! Furthermore: the Rivest and Vuillemin [19] result, which settles the ARK conjecture up to a constant factor, in fact confirms the Weak-EC for graph properties.

Theorem 1 (Rivest and Vuillemin). *If $\mathcal{P} : \{0, 1\}^{\binom{n}{2}} \rightarrow \{0, 1\}$ is a non-trivial monotone property of graphs on n vertices then $D(\mathcal{P}) = \Omega(n^2)$.*

It is interesting to note that the proof of equivalence in Kulkarni [7] does not hold between ARK and Weak-ARK. Hence: even though Weak-ARK is settled, the ARK is still wide open.

1.4 Our results on the Weak EC

In this paper we prove an analogue of Rivest and Vuillemin's result (Theorem 1) for 3-uniform hypergraphs. A property of 3-uniform hypergraphs on n vertices is a Boolean function $\mathcal{P} : \{0, 1\}^{\binom{n}{3}} \rightarrow \{0, 1\}$ whose variables are labeled by the $\binom{n}{3}$ potential edges of n -vertex 3-uniform hypergraphs and \mathcal{P} is invariant under relabeling of the vertices.

Theorem 2. *If $\mathcal{P} : \{0, 1\}^{\binom{n}{3}} \rightarrow \{0, 1\}$ is a non-trivial monotone property of 3-uniform hypergraphs on n vertices, then*

$$D(\mathcal{P}) = \Omega(n^3).$$

Our proof technique can be briefly described as follows: First we combine the combinatorial approach of Rivest and Vuillemin with the topological approach of Kahn, Saks, and Sturtevant to prove the result when $n = 3^k$. Then we use the 3^k case to prove the result for arbitrary n via an interesting application of

the famous Vinogradov's Theorem that asserts that every odd integer can be expressed as sum of three prime numbers.

Interestingly, we do not yet know how to generalize our proof technique to k -uniform hypergraphs. But in this context we are able to prove a partial result on 4-uniform hypergraphs.

Theorem 3. *Let $\mathcal{P} : \{0, 1\}^{n \times n \times n \times n} \rightarrow \{0, 1\}$ be a 4-uniform 4-partite hypergraph property of $4n$ -vertex hypergraphs. If \mathcal{P} is non-trivial and monotone, then*

$$D(\mathcal{P}) = \Omega(n^4).$$

The organisation of this paper is as follows. Section 2 contains the preliminaries. Section 3 contains the proof of $n = 3^k$ case. Section 4 uses 3^k case to prove the general case, in particular it contains the proof of Theorem 2. Section 5 contains some partial results for 4-uniform hypergraphs. Section 6 contains conclusion and open ends.

2 Preliminaries

In this paper $[n]$ denotes the set $\{1, \dots, n\}$.

2.1 Rivest-Vuillemin: combinatorial approach

In a beautiful paper, Rivest and Vuillemin show that the ARK Conjecture holds up to a constant factor, i.e., any non-trivial monotone graph property is weakly evasive. As an intermediate step [19] show the following:

Theorem 4 (Rivest-Vuillemin). *If n is a power of a prime number and $f : \{0, 1\}^n \rightarrow \{0, 1\}$ is any function invariant under a transitive permutation group such that $f(0, \dots, 0) \neq f(1, \dots, 1)$, then $D(f) = n$.*

In this paper we prove the weak-evasiveness of monotone properties of 3-uniform hyper-graphs, which extends the result of Rivest and Vuillemin for graph properties. Our proof is inspired by the one by Rivest and Vuillemin. Interestingly we use, in addition to the combinatorial approach of Rivest and Vuillemin, the powerful topological approach of Kahn, Saks, and Sturtevant combined with a deep theorem in number theory.

2.2 Kahn-Saks-Sturtevant: topological approach

In a seminal paper, Kahn, Saks, and Sturtevant [8] introduce a novel topological approach to settle the ARK Conjecture when the number of vertices of graphs is a power of a prime number. Their crucial observation was that non-evasiveness of monotone properties has a strong topological consequence, namely the corresponding simplicial complex is contractible to a point. Further they exploit this topological consequence via Oliver's Fixed Point Theorem [17] under the actions of certain special type of groups.

We say that a group Γ satisfies **Oliver's Condition** if there exist (not necessarily distinct) primes p, q such that Γ has a (not necessarily proper) chain of subgroups $\Gamma_2 \triangleleft \Gamma_1 \triangleleft \Gamma$ such that Γ_2 is a p -group, Γ_1/Γ_2 is cyclic, and Γ/Γ_1 is a q -group, where p -group means a group whose order is a power of a prime p .

Theorem 5 (Kahn-Saks-Sturtevant). *If Γ satisfies Oliver's Condition and acts transitively on the set S of variables, then for any non-trivial monotone Γ -invariant function $f : \{0, 1\}^S \rightarrow \{0, 1\}$, we have: $D(f) = |S|$.*

Kahn, Saks, and Sturtevant made the assumption that the number of vertices of the graph - n is a prime power and used the following group that satisfies Oliver's Condition:

$$AFF(n) := AGL(1, n),$$

the group of affine transformations $x \mapsto ax + b$ over the field \mathbb{F}_n of order n ; $a \in \mathbb{F}_n - \{0\}, b \in \mathbb{F}_n$. The two key properties of this group are that it is a cyclic extension of a p -group, i.e., it satisfies Oliver's Condition; moreover it acts doubly transitively on $[n]$, i.e., any (i, j) can be mapped to any (i', j') for $i \neq j$ and $i' \neq j'$.

In this paper, we make use of the $AFF(n)$, as well as another group-theoretic construction called wreath product. We recall the definition, and refer the readers to [20, Section 1.6] detailed discussion. For a finite set S , let $\text{Sym}(S)$ be the symmetric group on S . Let $G \leq \text{Sym}(S)$ and $H \leq \text{Sym}(T)$. The *wreath product* $G \wr H$ is a permutation group acting on $S \times T$, defined as follows. The *base group* of the wreath product is the direct product G^T , that is $|T|$ copies of G . For $t \in T$, the G_t independently acts on the corresponding copy $S \times \{t\}$. Specifically, for $(\omega, \delta) \in S \times T$, and $f \in G^T$, $(s, t)^f = (s^{f(t)}, t)$. $G \wr H$ also contains a subgroup H^* isomorphic to H , acting only on the second component of $S \times T$. That is for $h \in H$, $(s, t)^h = (s, t^h)$. $G \wr H$ is the group generated by G^S and H^* .

2.3 Prime-partition via Vinogradov's Theorem

The Goldbach Conjecture asserts that every even integer can be written as the sum of two primes. Vinogradov's Theorem [24] says that every sufficiently large odd integer m is the sum of three primes $m = p_1 + p_2 + p_3$. We use here Haselgroves version [5] of Vinogradov's theorem which states that we can require the primes to be roughly equal: $p_i \sim m/3$. This can be combined with the Prime Number Theorem to conclude that every sufficiently large even integer m is a sum of four roughly equal primes.

This fact was first used by Babai et. al. [1] to construct the group actions satisfying Oliver's Condition in order to show that any monotone property of sparse graphs is evasive.

3 3-uniform Hypergraphs: $n = 3^k$

We prove the following theorem in this section.

Theorem 6. *Let $n = 3^k$, and \mathcal{P} be 3-uniform hypergraph property of n -vertex hypergraphs. If \mathcal{P} is non-trivial and monotone, then $D(\mathcal{P}) = \Omega(n^3)$.*

Proof. Our proof strategy is inspired by the one by Rivest and Vuillemin's proof that non-trivial and monotone graph properties of graphs with $n = 2^k$ vertices are weakly evasive. The basic strategy is to set up a family of graphs $G_0 \subset G_1 \subset \dots \subset G_k$, among which there are two *adjacent* graphs G_ℓ and $G_{\ell+1}$ such that G_ℓ satisfies \mathcal{P} whereas $G_{\ell+1}$ does not. Now we start with the smaller graph G_ℓ and gradually *add* edges to it under the assumption that $D(\mathcal{P})$ is not $\Omega(n^3)$ and conclude that even after adding these edges the property \mathcal{P} is satisfied. Eventually, after adding sufficiently many edges this would lead to a contradiction as we would be able to conclude that $G_{\ell+1}$ satisfies the property.

Rivest and Vuillemin choose G_ℓ to be the disjoint union of $2^{n-\ell}$ cliques on 2^ℓ vertices. Further they use Theorem 4 to *add* the edges to finally lead to a contradiction. Similar to Rivest and Vuillemin, we start our proof by using Theorem 4 to add certain type of edges. However, while handling the 3-uniform hypergraph properties, we face more complications. The natural choice of G_i to be disjoint union of hyper-cliques seems to fail and Theorem 4 seems inadequate in dealing with all types of edges. We overcome this obstacle by suitably changing the family of graphs and by making use of the topological approach of Kahn, Saks, and Sturtevant (Theorem 5) to deal with the other type of edges.

3.1 Our choice of the graph family: cliques with spikes

To prove the theorem we consider the following family of hypergraphs on n vertices. For $j \in \{0, 1, \dots, k\}$, let G_j be the hypergraph defined as follows: firstly G_j contains a disjoint union of 3^{k-j} copies of cliques on 3^j vertices. Then if an edge $\{u, v, w\}$ satisfies that u, v are in the one clique while w is in another one, it is also included in G_j . We call such edges *spikes*.

As G_0 is the empty hypergraph, and G_k is the complete hypergraph, we see that G_0 satisfies \mathcal{P} while G_k does not as \mathcal{P} is non-trivial. This suggests that there exists $\ell \in \{0, 1, \dots, k-1\}$ such that G_ℓ satisfies the property while $G_{\ell+1}$ does not as \mathcal{P} is monotone.

Now collect the cliques in G_ℓ into three groups $V_1 \cup V_2 \cup V_3$, each group containing $3^{k-\ell-1}$ cliques. We then consider the property \mathcal{P}_1 induced by \mathcal{P} after fixing the values at the edges $\{\{u, v, w\} \mid u, v, w \in V_i, \text{ or } \{u, v, w\} \in G_\ell\}$ as in G_ℓ . Note that \mathcal{P}_1 is a non-trivial property, because \mathcal{P} is monotone and the graph $G_{\ell+1}$ is contained in the graph $G_\ell \cup E$ where E denotes the edges corresponding to the domain of \mathcal{P}_1 .

3.2 Two types of edges

The edges not fixed in \mathcal{P}_1 are of two types:

Type 1 $T_1 = \{\{v_1, v_2, v_3\} \mid v_i \in V_i, i \in [3]\}$.

Type 2 $T_2 = \{\{u, v, w\} \mid u, v \in V_i, w \in V_j, i \neq j\}$. Note that v, w cannot come from the same clique otherwise it would have been fixed.

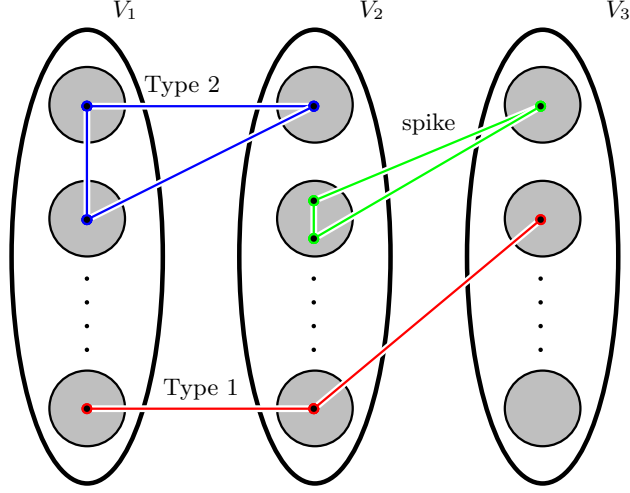


Fig. 1. An illustration of the spike (green), and Type 1 (red) and Type 2 (blue) edges.

Before going on we define two group actions on V_1 . Firstly, $H_1 = \mathbb{Z}_{3^\ell} \wr \mathbb{Z}_{3^{k-1-\ell}}$ acts on V_1 , where $3^{k-1-\ell}$ copies of \mathbb{Z}_{3^ℓ} act independently on the $3^{k-1-\ell}$ cliques, and $\mathbb{Z}_{3^{k-1-\ell}}$ permutes among the cliques. Secondly, we define the group action of $H_2 = \mathbb{Z}_{3^\ell} \wr \text{AFF}(3^{k-1-\ell})$ on V_1 similarly to H_1 . That is, $3^{k-1-\ell}$ copies of \mathbb{Z}_{3^ℓ} act independently on the $3^{k-1-\ell}$ cliques, and $\text{AFF}(3^{k-1-\ell})$ acts on the cliques in doubly-transitive way. Recall that for a vector space V , $\text{AFF}(V)$ is the affine group on V . H_1 and H_2 are both subgroups of the automorphism group of induced subgraph of G_ℓ on V_1 . Then note that H_1 is a 3-group, and H_2 is transitive on $\{\{u, v\} \mid u \text{ and } v \text{ are from different cliques}\}$. Finally, it can be verified that H_2 belongs to the group class as in Theorem 5.

Table 1. Groups used to handle Type 1 and Type 2 edges

Type 1	$(\mathbb{Z}_{3^\ell} \wr \mathbb{Z}_{3^{k-1-\ell}}) \times (\mathbb{Z}_{3^\ell} \wr \mathbb{Z}_{3^{k-1-\ell}}) \times (\mathbb{Z}_{3^\ell} \wr \mathbb{Z}_{3^{k-1-\ell}})$
Type 2	$(\mathbb{Z}_{3^\ell} \wr \text{AFF}(3^{k-1-\ell})) \times (\mathbb{Z}_{3^\ell} \wr \mathbb{Z}_{3^{k-1-\ell}})$

3.3 Adding Type 1 edges

Now we consider the property \mathcal{P}_2 induced by \mathcal{P}_1 by setting Type 2 edges to be absent. Note that the number of Type 1 edges is $3^{3(k-1)}$, thus a prime power. Let $H_1 \times H_1 \times H_1$ act on $V_1 \times V_2 \times V_3$ in a natural way: each copy of H_1 acts on vertices of V_i independently. It is seen that this action preserves the fixed subgraph, and \mathcal{P}_2 is invariant under this action. If after adding all Type 1 edges \mathcal{P}_2 would *not* be satisfied, then by the Rivest-Vuillemin theorem, \mathcal{P}_2 is evasive. That is $D(\mathcal{P}) \geq D(\mathcal{P}_2) = 3^{3(k-1)} = \Omega(n^3)$ and we would be done.

3.4 Adding Type 2 edges

Let \mathcal{P}_3 be the property induced by \mathcal{P}_1 by setting Type 1 edges to be present. The discussion from last paragraph suggests that \mathcal{P}_3 is a non-trivial property, and note that \mathcal{P}_3 only has Type 2 edges left unfixed. For $i, j \in [3], i \neq j$, let $T_2(i, j) = \{\{u, v, w\} \mid u, v \in V_i, w \in V_j\}$. Let \mathcal{P}_4 be the property induced by \mathcal{P}_3 by setting edges in $T_2 \setminus T_2(1, 2)$ to be absent. Note that $|T_2(1, 2)| = \Omega(n^3)$.

(Oliver's Condition Holds) Consider the group $H = H_2 \times H_1$ acting on $V_1 \times V_2$ in a natural way. It is verified that H preserves the structure of the fixed graph, and \mathcal{P}_3 is invariant under H . It is easy to check that H belongs to the group class described in Theorem 5. This allows us to apply Theorem 5 (the Kahn, Saks and Sturtevant Theorem) to conclude that either \mathcal{P}_4 is trivial; if not then its query complexity is $\Omega(n^3)$ and we would be done.

(Orbits are large) Here we use a key property of the action of H , namely that the orbit of any edge is of large size: $\Omega(n^3)$. Thus we can add edges in $T_2(1, 2)$ to get another restriction \mathcal{P}_5 . Then we use the same group as above for $V_2 \times V_3$ to add edges in $T_2(2, 3)$.

3.5 Deriving a contradiction

Continuing this way we can keep adding $T_2(i, j)$ edges while maintaining that the hyper-graph still satisfies the property. But then we would get $G_{\ell+1}$ as a subgraph which by our choice of ℓ , does not satisfy the property. Contradiction!

4 3-uniform Hypergraphs: General n

In this section we prove the main theorem.

Theorem 1.5, restated. If $\mathcal{P} : \{0, 1\}^{\binom{n}{3}} \rightarrow \{0, 1\}$ is a non-trivial monotone property of 3-uniform hypergraphs on n vertices, then $D(\mathcal{P}) = \Omega(n^3)$.

Proof. The natural way of extending Rivest and Vuillemin's argument for 3-uniform hyper graphs for arbitrary n leads to analysis of several types of edges. We do not know an easy way to handle this via combinatorial approach. We can use the topological approach together with an interesting theorem about partitioning an integer into prime numbers to patch up the 3^k case to arbitrary n .

4.1 Prime-partition of n via Vinogradov's Theorem

We distinguish two cases: (Case 1) n is even and (Case 2) n is odd. Let us consider Case 1: n is even. The other case can be handled in a similar fashion. Let k be the largest power of 3 that does not exceed n . Since n is even, we can write (using the above mentioned Hasegrove's version of Vinogradov's Theorem) $n = p_1 + p_2 + p_3 + 3^{k-1}$, where p_i s are prime numbers and $p_i \sim p_j$. Moreover: note that by our choice of k we can assume: $p_i \leq 3^k$.

4.2 Patching up 3^k case to general n

We partition $[n]$ into parts of size p_1, p_2, p_3 and 3^{k-1} as described in the previous section. Let \mathcal{P} be a non-trivial monotone property of 3-uniform hyper-graphs on n vertices. Theorem 6 allows us to conclude that either (a) $D(\mathcal{P}) = \Omega(n^3)$ or (b) *any* 3^k vertex (hyper) clique satisfies \mathcal{P} . In Case (a) we are done. So let us assume that we are in Case (b). Since $p_1 \leq 3^k$ and since \mathcal{P} is monotone, we may assume that the clique on p_1 vertices satisfies the property. Now we assume that the clique on p_1 vertices is present and restrict our attention to the induced property \mathcal{P}_2 of 3-uniform hyper-graphs on $p_2 + p_3 + 3^{k-1}$ vertices. Again using the fact that $p_2 \leq 3^k$ we can assume that the clique on p_2 vertices is also present in addition to the clique on p_1 vertices. Now we move our attention to the induced property \mathcal{P}_2 on $p_3 + 3^{k-1}$ vertex graphs. In one more step, we can move our attention to the induced property \mathcal{P}_3 on 3^{k-1} graphs which assumes that the cliques on the p_1, p_2 , and p_3 vertices are present. Finally, with the use of Theorem 6, we can conclude that the clique on the 3^{k-1} vertices is also present; if not then we could already conclude $D(\mathcal{P}) = \Omega(n^3)$.

4.3 Two types of edges

Now we have a restriction \mathcal{P}' of our original property \mathcal{P} in which the cliques on p_1, p_2, p_3 and 3^{k-1} vertices are present. We partition the absent edges into two types:

- Type A** the three endpoints of the edges belong to different cliques;
- Type B** two of the three endpoints belong to one clique and the remaining endpoint belongs to a different clique.

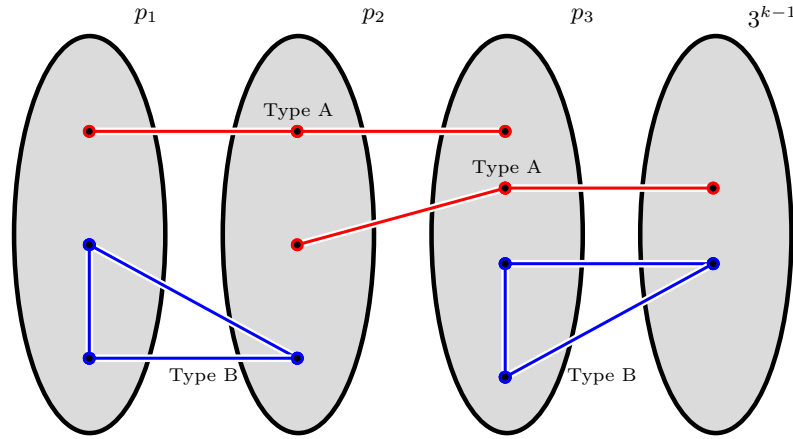


Fig. 2. An illustration of Type A (red) and Type B (blue) edges.

Table 2. Groups used to handle Type A and Type B edges

Type A	$\mathbb{Z}_{p_1} \times \mathbb{Z}_{p_2} \times \mathbb{Z}_{p_3}$, and $\mathbb{Z}_{p_i} \times \mathbb{Z}_{p_j} \times \mathbb{Z}_{3^{k-1}}$
Type B	$AF\!F(p_i) \times \mathbb{Z}_{p_j}$, and $AF\!F(p_i) \times \mathbb{Z}_{3^{k-1}}$, and $AF\!F(3^{k-1}) \times \mathbb{Z}_{p_i}$

4.4 Adding Type A edges

Firstly: we conclude that all Type A edges must also be present; if not then $D(\mathcal{P}) = \Omega(n^3)$. For this we use the following two types of groups: $\mathbb{Z}_{p_1} \times \mathbb{Z}_{p_2} \times \mathbb{Z}_{p_3}$ and $\mathbb{Z}_{p_i} \times \mathbb{Z}_{p_j} \times \mathbb{Z}_{3^{k-1}}$.

4.5 Adding Type B edges

Secondly: after adding all Type A edges we can conclude that all Type B edges must also be present; if not then $D(\mathcal{P}) = \Omega(n^3)$. For this we use the following three types of groups: $AF\!F(p_i) \times \mathbb{Z}_{p_j}$ and $AF\!F(p_i) \times \mathbb{Z}_{3^{k-1}}$ and $AF\!F(3^{k-1}) \times \mathbb{Z}_{p_i}$.

(Oliver's Condition Holds) It is easy to check that all the groups that we use are the right ones for using the topological approach, i.e., they are “ q -group extension of cyclic extension of p -groups,” i.e., they satisfy Oliver's Condition.

(Orbits are large) A crucial property that we used in our proof is that the orbit of any edge under any of our group actions is large: $\Omega(n^3)$.

4.6 Deriving a contradiction

After adding both Type A and Type B edges to the cliques on the p_1, p_2, p_3 and 3^{k-1} vertices, we can conclude that the clique on n vertices must satisfy \mathcal{P} ; this contradicts with our initial assumption that \mathcal{P} is non-trivial.

This completes the proof of Theorem 2.

5 4-uniform 4-partite Hypergraphs

In this section we prove the weak evasiveness for properties of 4-uniform 4-partite hypergraphs. **Theorem 1.6, restated.** Let $\mathcal{P} : \{0, 1\}^{n \times n \times n \times n} \rightarrow \{0, 1\}$ be a 4-uniform 4-partite hypergraph property of $4n$ -vertex hypergraphs. If \mathcal{P} is non-trivial and monotone, then $D(\mathcal{P}) = \Omega(n^4)$.

Proof: If n is prime, the result directly follows from Theorem 4.

In the case when n is not prime, let p be a prime number such that $p < n < 2p$. Let $V = V_1 \cup V_2 \cup V_3 \cup V_4$, $|V_i| = n$ be the vertex set. The strategy is again by contrapositive: assume $D(\mathcal{P})$ is not of $\Omega(n^4)$. Then we shall start from the empty graph, and then add the edges with different types while keeping the value of the property not change. Finally we will get that the complete graph satisfies the property, which contradicts to the condition of being a non-trivial property.

Let G_0 be the empty graph; thus $f(G_0) = 0$. Let $V_i = A_i \cup B_i$, where A_i is a vertex set of size p ($i = 1, 2, 3, 4$), and $B_i = V_i \setminus A_i$.

5.1 Adding edges in $A_1 \times A_2 \times A_3 \times A_4$

Consider a restriction \mathcal{P}_1 of \mathcal{P} where all the variables outside $A_1 \times A_2 \times A_3 \times A_4$ are set to be 0. \mathcal{P}_1 is a monotone transitive invariant function with p^4 variables, by Theorem 4 \mathcal{P}_1 is trivial, otherwise $D(\mathcal{P}_1) = p^4 = \Omega(n^4)$. Let G_1 be the graph with edges $A_1 \times A_2 \times A_3 \times A_4$, thus $f(G_1) = 0$.

5.2 Adding edges in $B_1 \times A_2 \times A_3 \times A_4$

We will add all the edges in $B_1 \times A_2 \times A_3 \times A_4$ to G_1 , resulting in a graph G_2 with edges $V_1 \times A_2 \times A_3 \times A_4$. Before doing that we consider a graph G'_1 with edges $B_1 \times A_2 \times A_3 \times A_4$. Since $p > n - p$, from the monotone and symmetry condition, we have $f(G'_1) \leq f(G_1) = 0$. Consider a restriction \mathcal{P}_2 of \mathcal{P} where all the edges in G'_1 are set to 1 and all the edges outside $V_1 \times A_2 \times A_3 \times A_4$ are set to 0. It is clear that \mathcal{P}_2 is a monotone transitive invariant function with p^4 variables, thus from Theorem 4 \mathcal{P}_2 is a constant, otherwise $D(\mathcal{P}_2) = p^4 = \Omega(n^4)$. Hence we get $f(G_2) = 0$.

5.3 Adding edges in $V_1 \times B_2 \times A_3 \times A_4$

Similar to the previous step, we first use the monotone and symmetry condition to “delete” some edges from G_2 . Let G'_2 be the graph with edges $V_1 \times B_2 \times A_3 \times A_4$. From the monotone and symmetry condition, $f(G'_2) \leq f(G_2) = 0$. Consider the restriction \mathcal{P}_3 of \mathcal{P} where all the edges in G'_2 are set to 1 and all the edges outside $V_1 \times V_2 \times A_3 \times A_4$ are set to 0. It is easy to see that \mathcal{P}_3 can be further partitioned into two properties isomorphism to \mathcal{P}_1 and \mathcal{P}_2 , respectively. By repeating the steps in Section 5.1 and 5.2 we conclude that \mathcal{P}_3 is trivial, otherwise $D(\mathcal{P}_3) = \Omega(n^4)$. Hence $f(G_3) = 0$, where G_3 is the graph with edges $V_1 \times V_2 \times A_3 \times A_4$.

Adding edges in $V_1 \times V_2 \times B_3 \times A_4$ and $V_1 \times V_2 \times V_3 \times B_4$

These two steps are similar to the previous step, and we omit them here. After doing these steps, we get that the value of the complete graph is also 0, which contradicts the non-trivial condition. \square

Remark 1. We note that the proof strategy for Theorem 1.6 can be extended to show that any k -uniform k -partite hypergraph property is weakly evasive, when k is a constant. On the other hand, we do not know how to prove the weak evasiveness for 4-uniform hypergraph properties.

6 Conclusion

In this paper we are able to confirm a special case of the Weak-EC. In particular, we have shown that any non-trivial monotone property of 3-uniform hypergraphs is weakly evasive. It is interesting to see how far can one generalize our results.

Question 1. Is any non-trivial monotone property of k -uniform hypergraphs weakly evasive?

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